# Towards Service Differentiation on the Internet

#### from

"New Internet and Networking Technologies for Grids and High-Performance Computing", tutorial given at IEEE HOTI 2006, Stanford, California August 25th, 2006

C. Pham
University of Pau, France
LIUPPA laboratory

## Revisiting the same service for all paradigm



No delivery guarantee

**INTERNET** 



#### Enhancing the best-effort service



**Introduce Service Differentiation** 



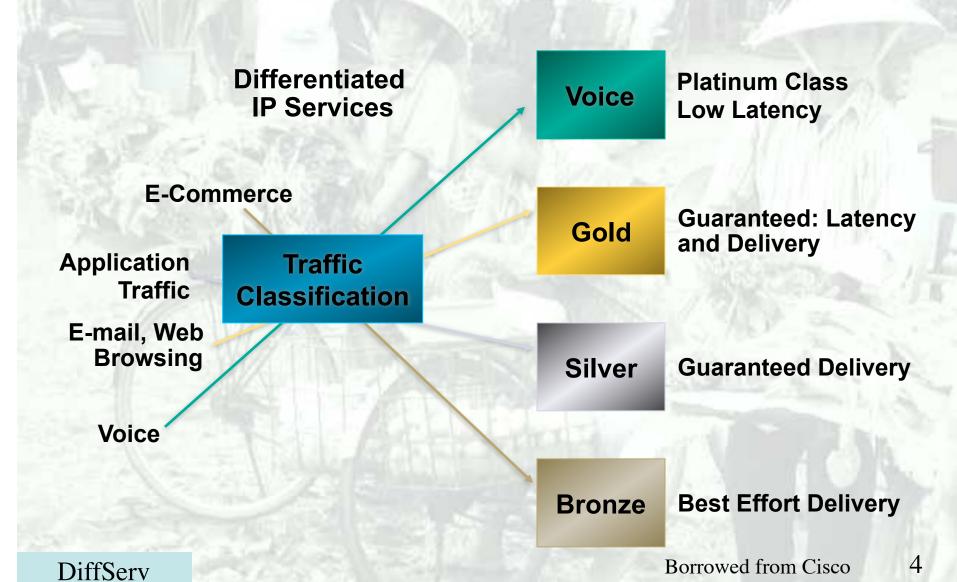
#### Service Differentiation

The real question is to choose which packets shall be dropped. The first definition of differential service is something like "not mine." -- Christian Huitema

- □ Differentiated services provide a way to specify the relative priority of packets
- □ Some data is more important than other
- People who pay for better service get it!



#### Divide traffic into classes



### Design Goals/Challenges

- □ Ability to charge differently for different services
- ■No per flow state or per flow signaling
- ☐ All policy decisions made at network boundaries
  - Boundary routers implement policy decisions by tagging packets with appropriate priority tag
- □ Traffic policing at network boundaries
- Deploy incrementally: build simple system at first, expand if needed in future

IP implementation: DiffServ

Flow 1

Flow:

Flow 4

No per flow state in the core

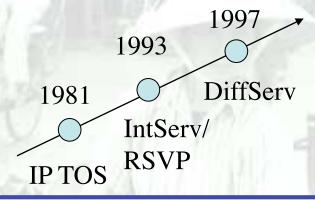
IP packet

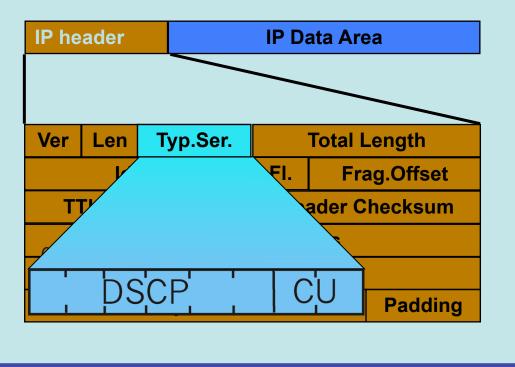


10Gbps=2.4Mpps with 512-byte packets

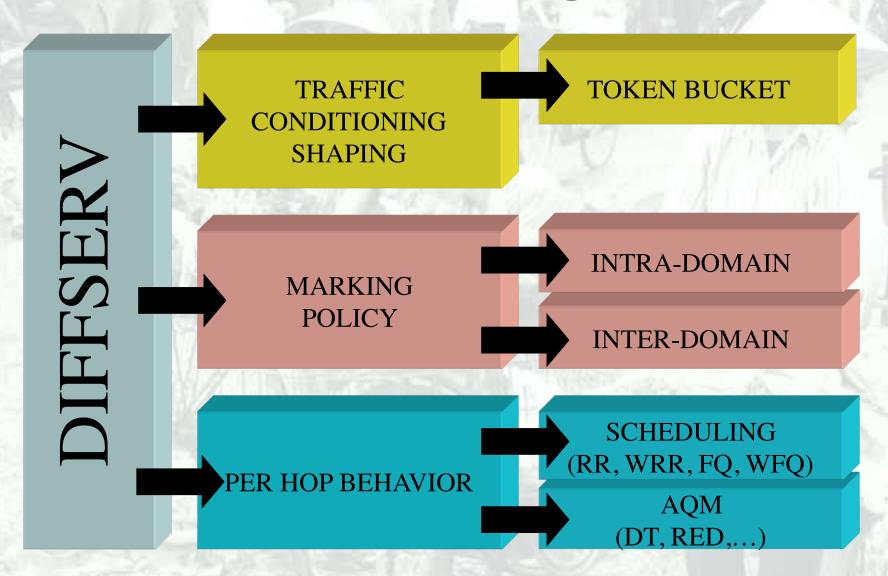
Stateful approaches scalable at gigabit rates

6 bits used for Differentiated Service Code Point (DSCP) and determine PHB that the packet will receive RFC 2475





# DiffServ building blocks



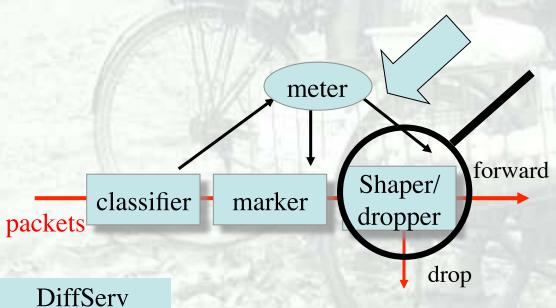
### Traffic Conditioning

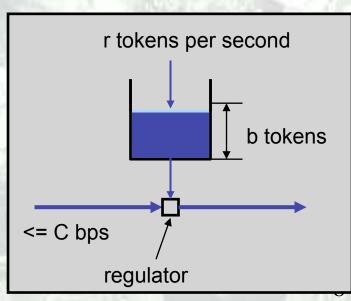
User declares traffic profile (eg, rate and burst size); traffic is metered and shaped if non-conforming

5Mbps

SLA 2Mbps

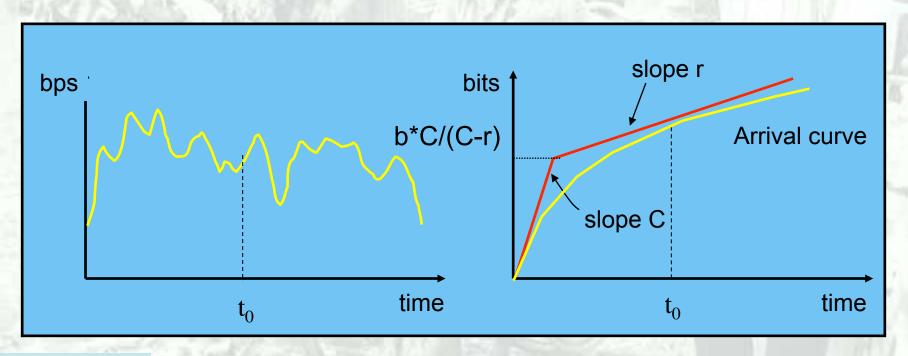
Service



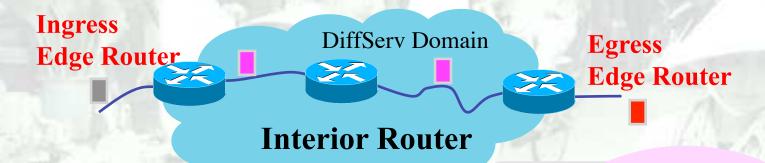


# Token Bucket for traffic characterization

□Given b=bucket size, C=link capacity and r=token generation rate



#### Differentiated Architecture

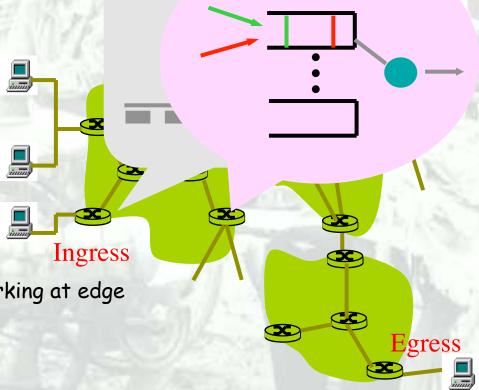


#### Marking:

per-flow traffic management marks packets as in-profile and outprofile

#### Per-Hop-Behavior (PHB):

per class traffic management Ingress buffering and scheduling based on marking at edge preference given to in-profile packets



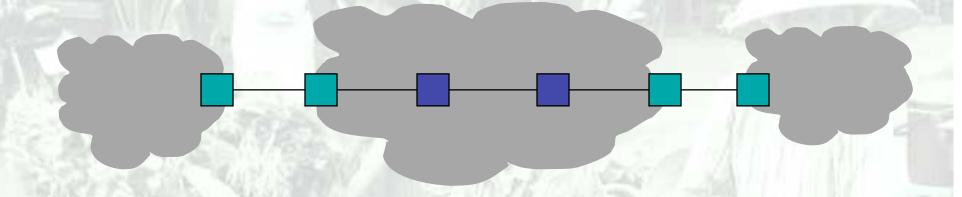
scheduling

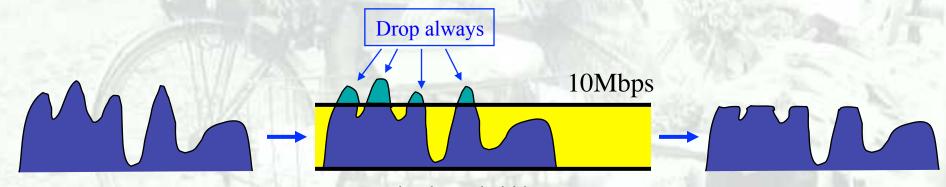
#### Pre-defined PHB

- □ Expedited Forwarding (EF, premium):
  - departure rate of packets from a class equals or exceeds a specified rate (logical link with a minimum guaranteed rate)
  - □ Emulates leased-line behavior

- □ Assured Forwarding (AF):
  - □ 4 classes, each guaranteed a minimum amount of bandwidth and buffering; each with three drop preference partitions
  - Emulates frame-relay behavior

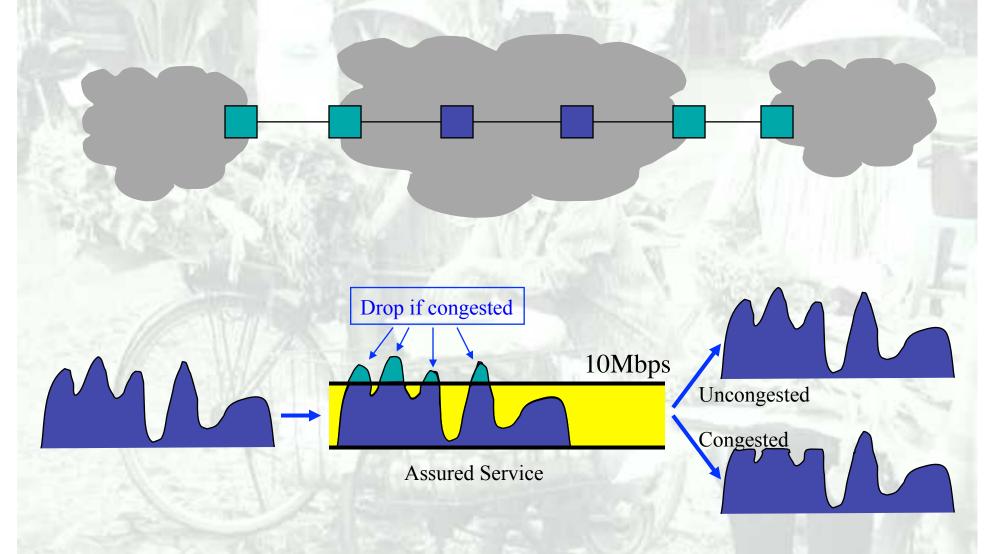
# Premium Service Example



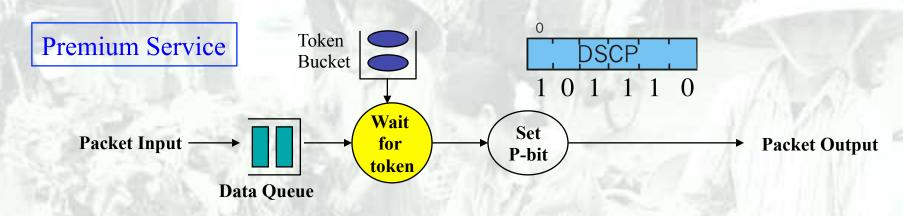


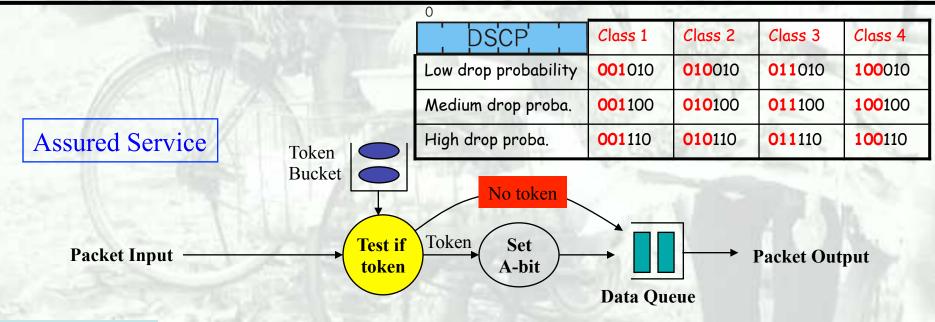
Fixed Bandwidth

# Assured Service Example

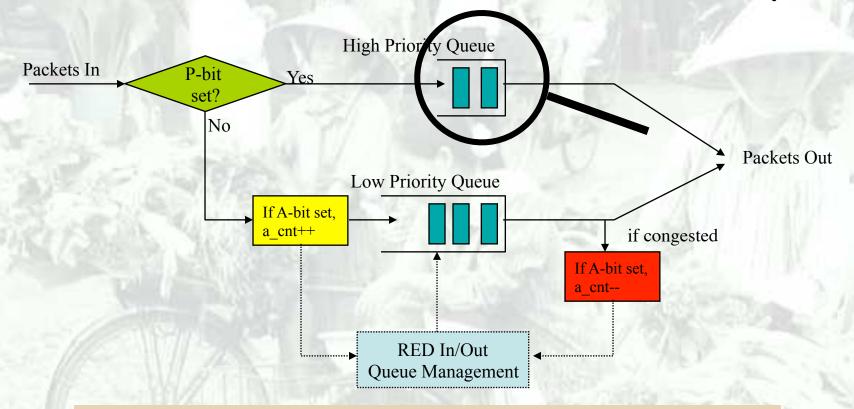


### Border Router Functionality





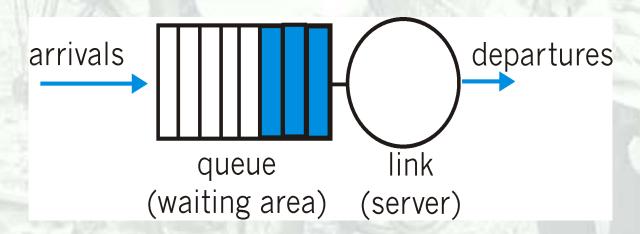
#### Internal Router Functionality

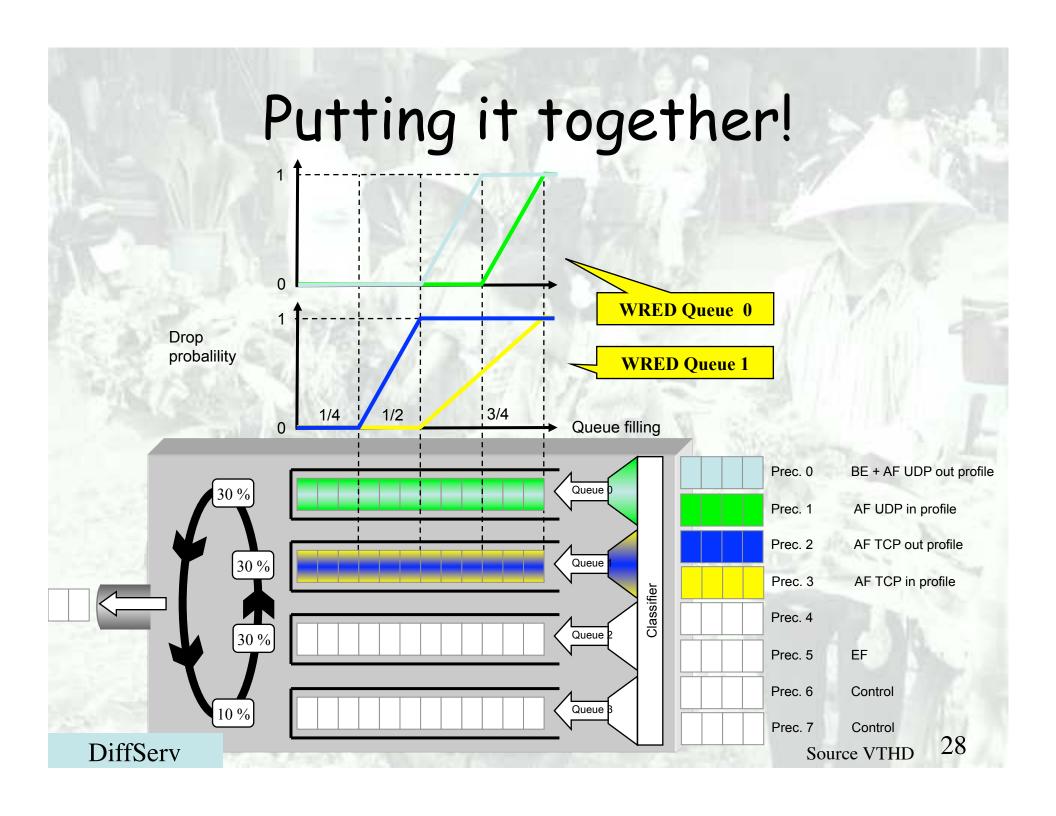


A DSCP codes aggregates, not individual flows
No state in the core
Should scale to millions of flows

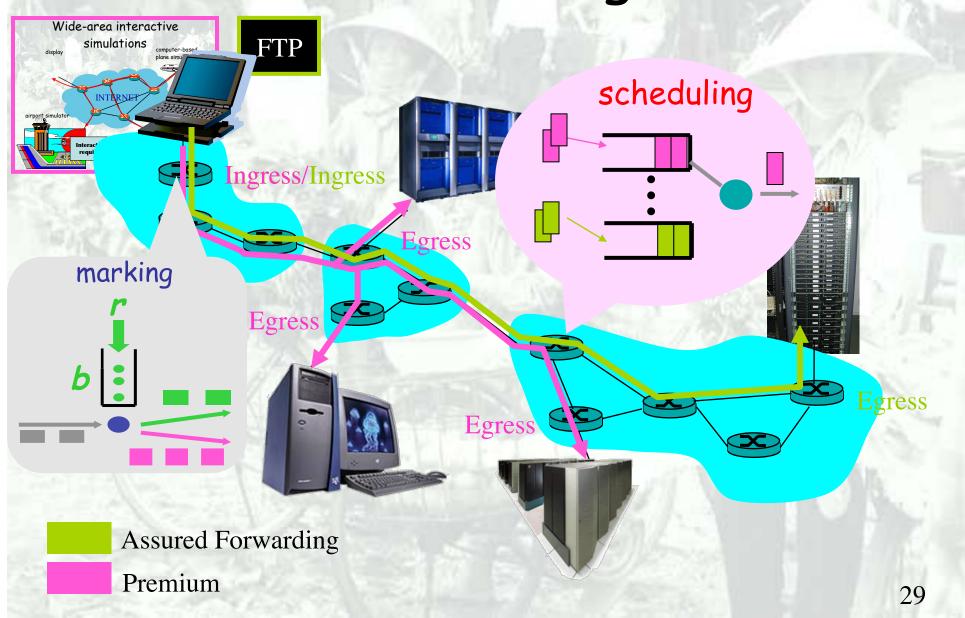
### Scheduling

- □ DiffServ PHB relies mainly on scheduling
  - choose the next packet for transmission
  - □ FIFO: in order of arrival to the queue; packets that arrive to a full buffer are either discarded, or a discard policy is defined.
  - ☐ More complex policies: FCFS, PRIORITY, EDD...

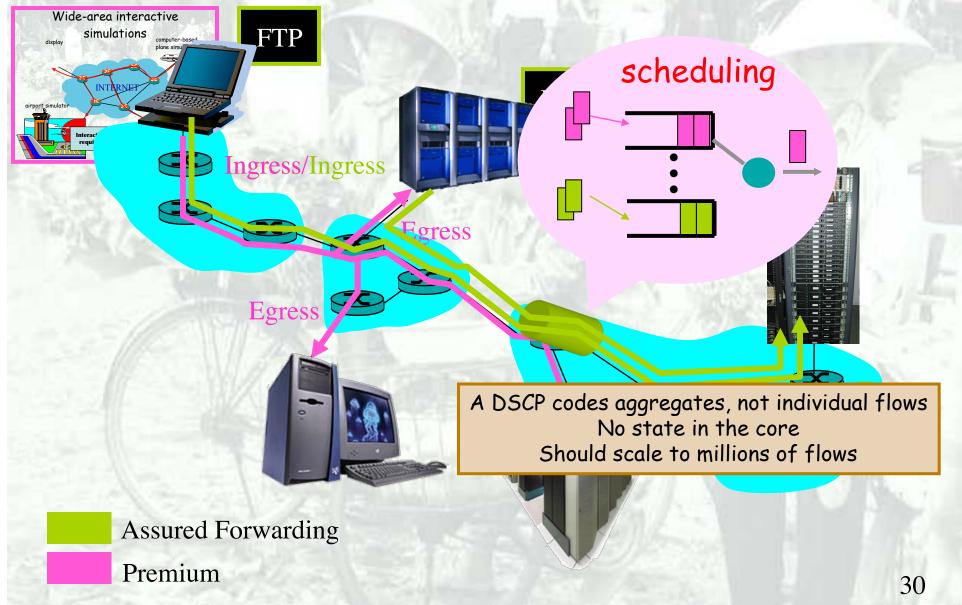




# DiffServ for grids

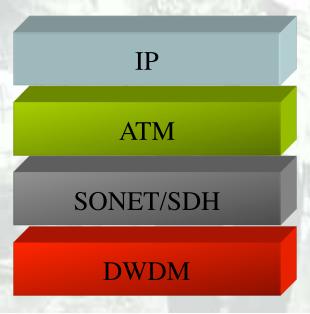


# DiffServ for grids (con't)

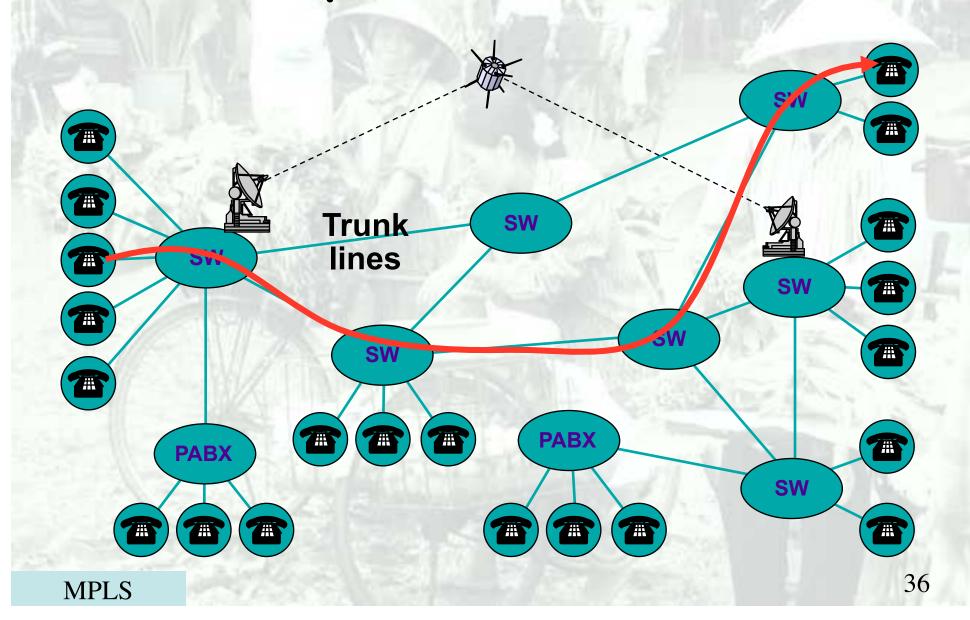


### Bandwidth provisioning

- □ DWDM-based optical fibers have made bandwidth very cheap in the backbone
- On the other hand, dynamic provisioning is difficult because of the complexity of the network control plane:
  - □ Distinct technologies
  - Many protocols layers
  - Many control software



# The telephone circuit view



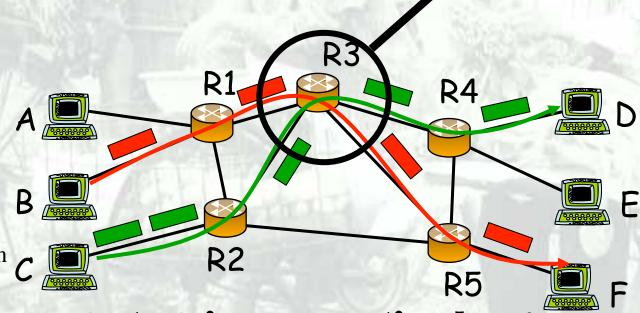
## Advantages of circuits

- Provides the same path for information of the same connection: less out-of-order delivery
- □ Easier provisioning/reservation of network's resources: planning and management features

#### Back to virtual circuits

□Virtual circuit refers to a connection oriented network/link layer: e.g. X.25, Frame Relay, ATM

Virtual
Circuit
Switching:
a path is defined
for each connection



**But IP is connectionless!** 

# Why virtual circuit?

□Initially to speed up router's forwarding tasks: X.25, Frame Relay, ATM.

We're fast

enough!

Now: Virtual circuits for traffic engineering!

#### Virtual circuits in IP networks

- Multi-Protocol Label Switching
  - □ Fast: use label switching → LSR
  - Multi-Protocol: above link layer, below network layer
  - □ Facilitate traffic engineering

PPP Header(Packet over SONET/SDH)

**Ethernet** 

**Frame Relay** 

**PPP Header MPLS** Header

**Ethernet Hdr MPLS Header** 

**FR Hdr** 

**MPLS** Header

**Layer 3 Header** 

IP

**MPLS** 

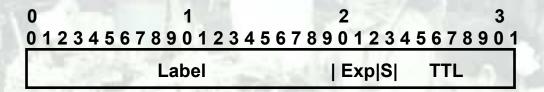
LINK

Layer 3 Header

**Layer 3 Header** 

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#### Label structure



Label = 20 bits

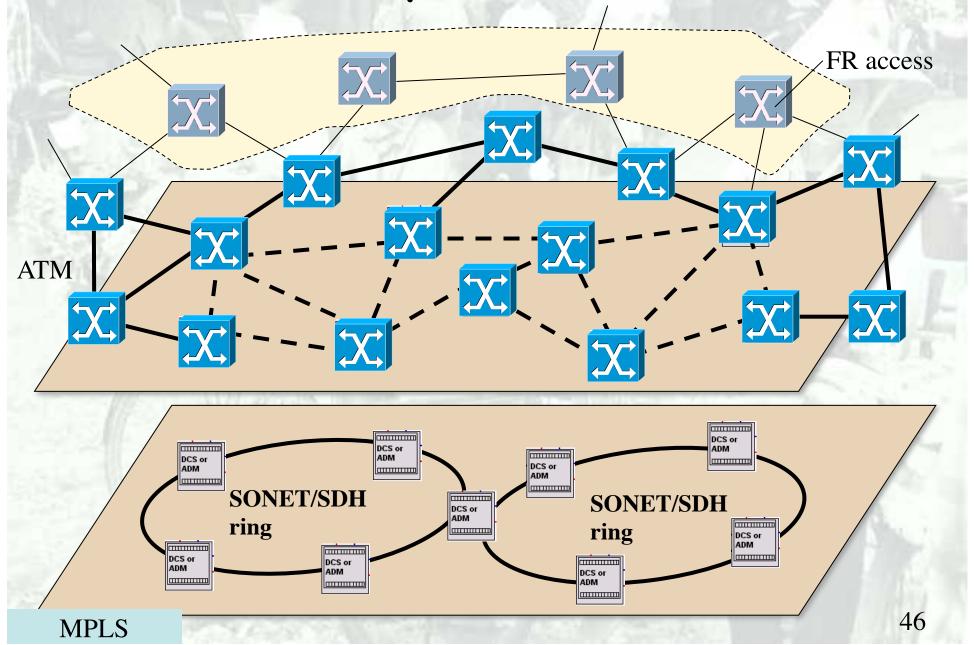
Exp = Experimental, 3 bits

S = Bottom of stack, 1bit

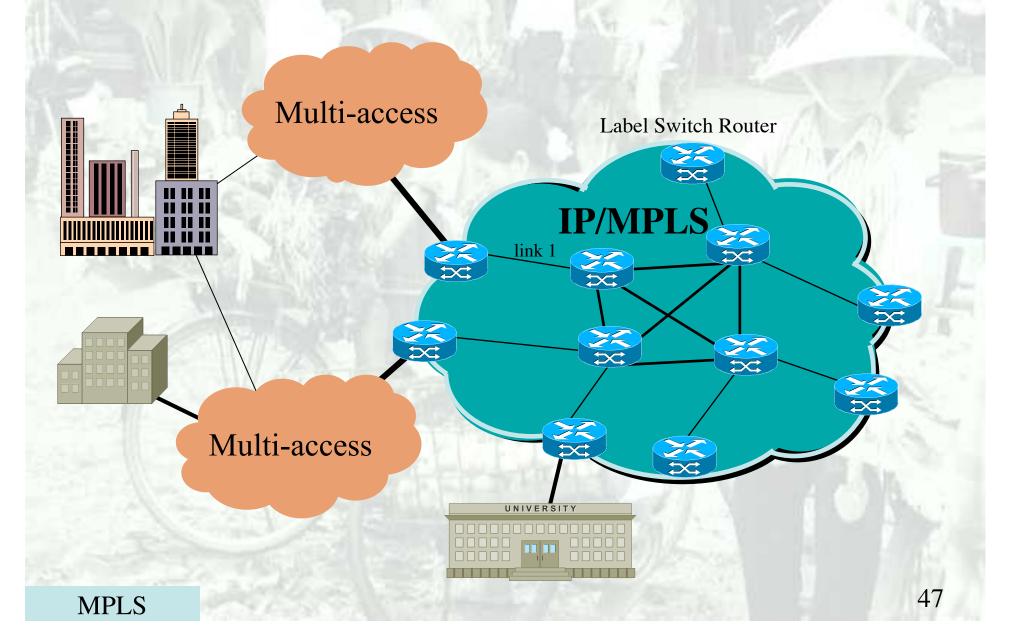
TTL = Time to live, 8 bits

- More than one label is allowed -> Label Stack
- MPLS LSRs always forward packets based on the value of the label at the top of the stack

# From multilayer networks...



#### ... to IP/MPLS networks

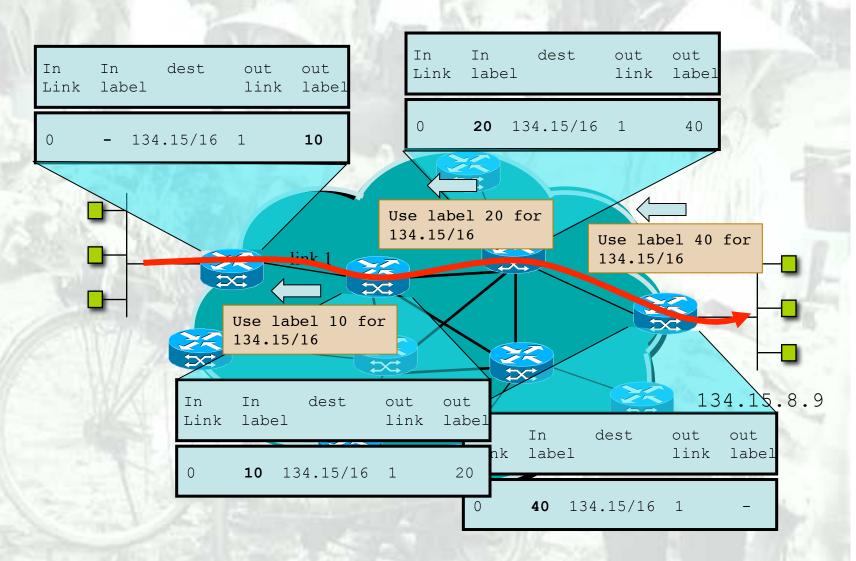


#### MPLS operation

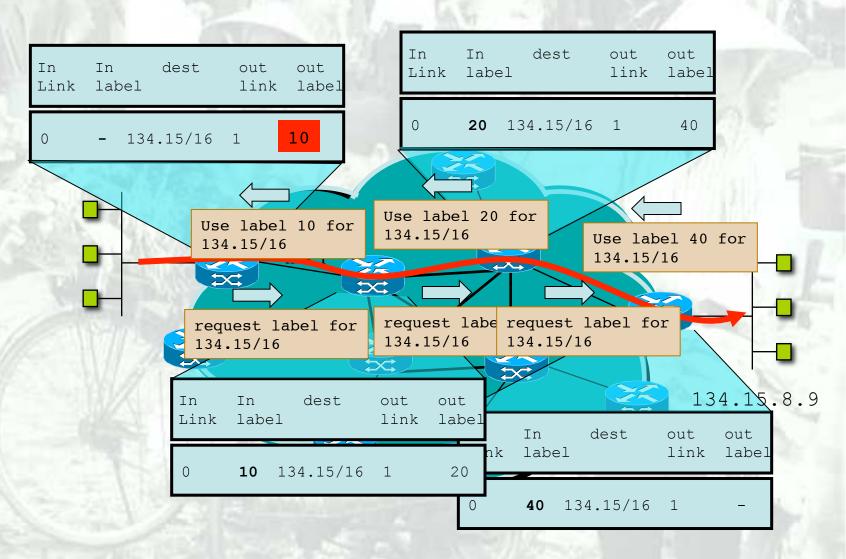
4. LSR at egress 1a. Routing protocols (e.g. OSPF-TE, IS-IS-TE) removes label and exchange reachability to destination networks delivers packet 1b. Label Distribution Protocol (LDP) Label Switch Router establishes label mappings to destination network dest out src 134.15.8.9 134.15/16 1/10 140.134/16 1/26 2. Ingress LSR receives packet and "label"s packets 3. LSR forwards Source Yi Lin, modified C. Pham packets using label switching

**MPLS** 

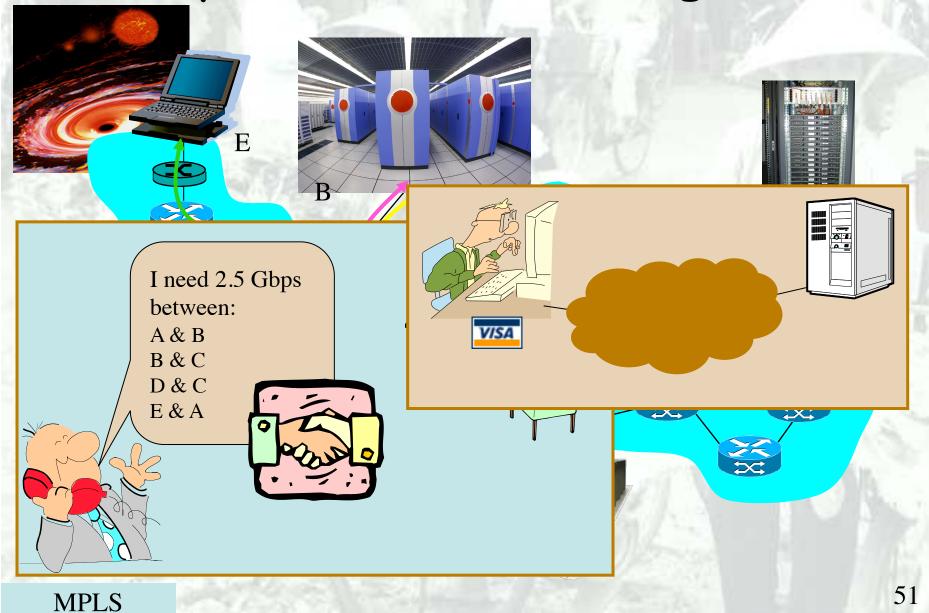
#### Label Distribution



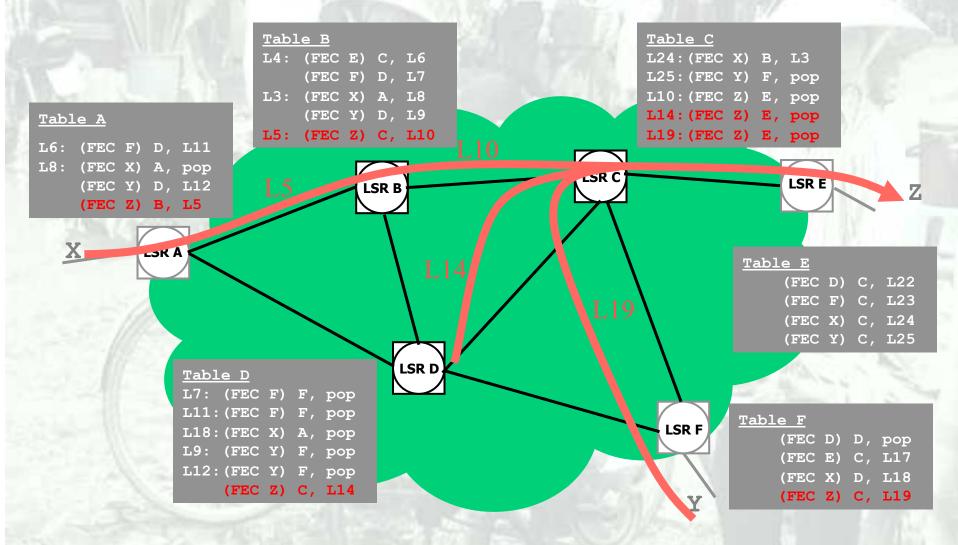
#### Label Distribution (con't)



# Dynamic circuits for grids



# Forwarding Equivalent Class: high-level forwarding criteria



# Forwarding Equivalent Class

A FEC aggregates a number of individual flows with the same characteristics: IP prefix, router ID, delay or bandwidth constraints...

B, L3 F, pop

#### **One possible utilization of FEC** Table A L6: (FEC F) **FEC A** (FEC X) (FEC Y) **L34** (FEC Z) **FTP** FEC B **Application FEC** L45 **Traffic** Classification E-mail FEC C Web **L07 Browsing Ingress Voice** LSR

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C, L22

C, L23

C, L25

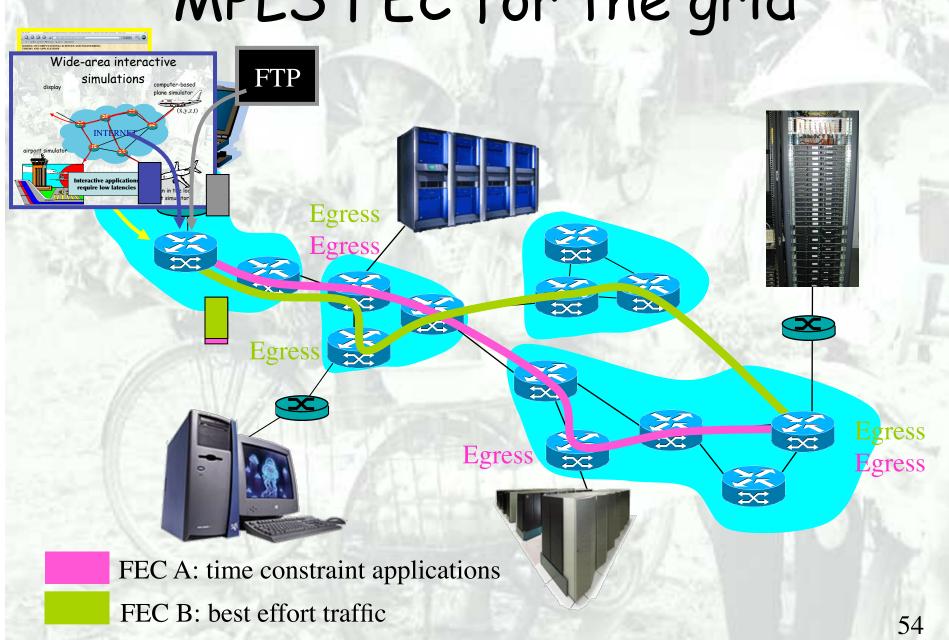
L24

pop

**L17** 

L18





#### Label & FEC

- ☐ Independent LSP control
  - ☐ An LSR binds a label to a FEC, whether or not the LSR has received a label from the next-hop for the FEC
  - ☐ The LSR then advertises the label to its neighbor
- Ordered LSP control
  - ☐ An LSR only binds and advertises a label for a particular FEC if:
    - it is the egress LSR for that FEC or
    - · it has already received a label binding from its next-hop

#### Label Distribution Protocols

- **ULDP** 
  - · Maps unicast IP destinations into labels
- RSVP-TE, CR-LDP
  - · Used in traffic engineering
- □ BGP
  - · External labels (VPN)
- ☐ PIM
  - · For multicast states label mapping

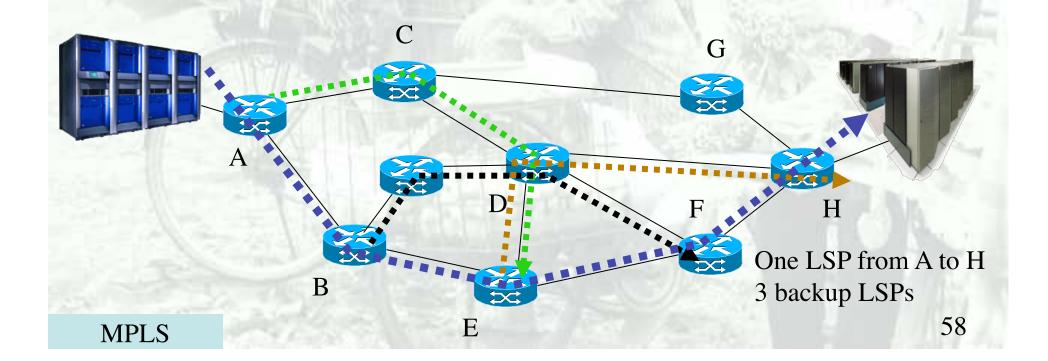
### MPLS for resiliency

MPLS FastReroute

- □Intended to provide SONET/SDH-like healing capabilities
- Selects an alternate route in tenth of ms, provides path protection
- □ Traditional routing protocols need minutes to converge!
- □ FastReroute is performed by maintaining backup LSPs

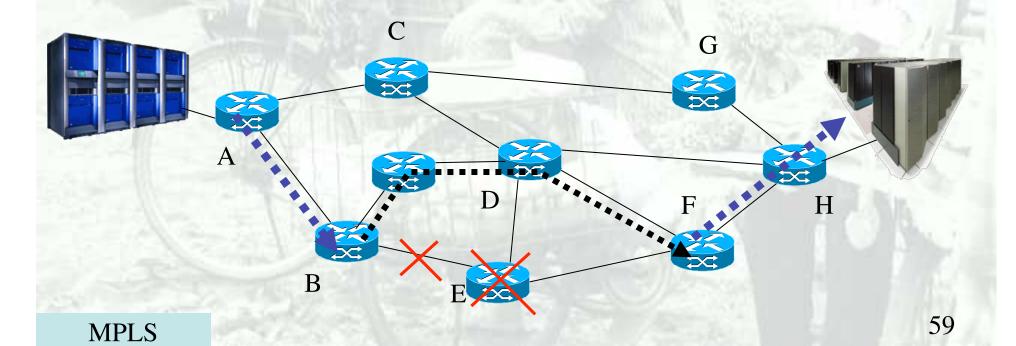
# MPLS for resiliency, con't Backup LSPs

- ☐ One-to-one
- Many-to-one: more efficient but needs more configurations



# MPLS for resiliency, con't Recovery on failures

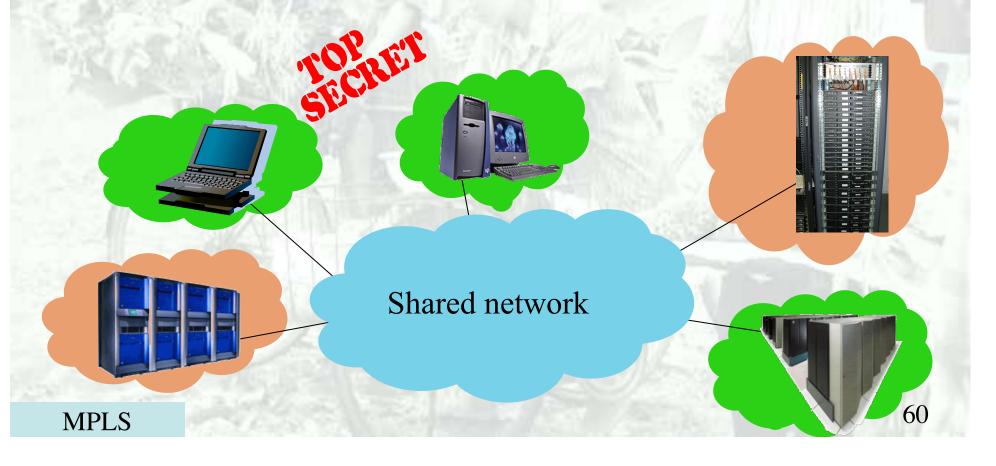
- □ Suppose E or link B-E is down...
- □B uses detour around E with backup LSP



#### MPLS for VPN

(Virtual Private Networks)

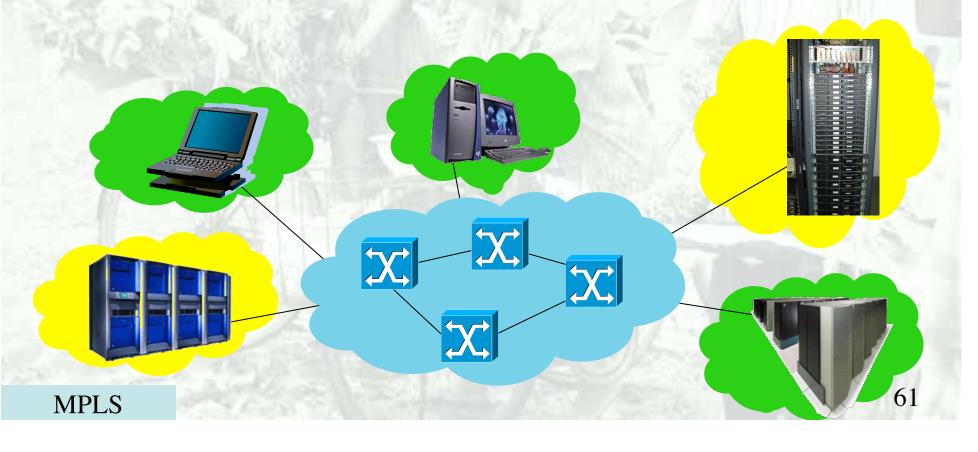
□ Virtual Private Networks: build a secure, confidential communication on a public network infrastructure using routing, encryption technologies and controlled accesses



### MPLS for VPN, con't

The traditional way of VPN

□Uses leases lines, Frame Relay/ATM infrastructures...

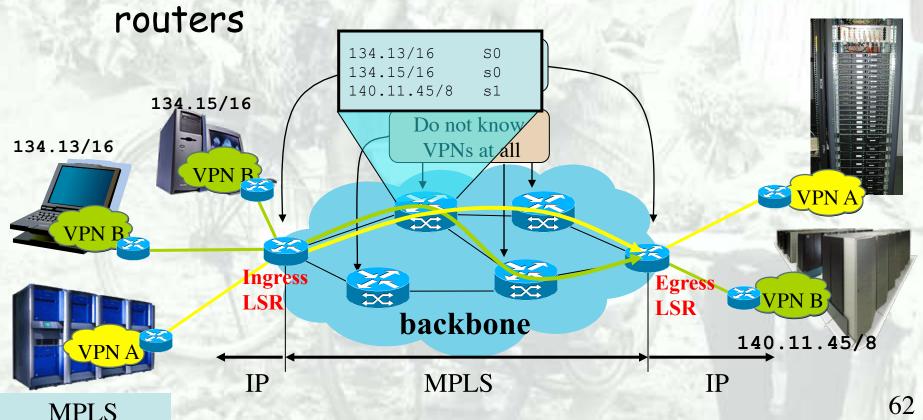


#### MPLS for VPN, con't

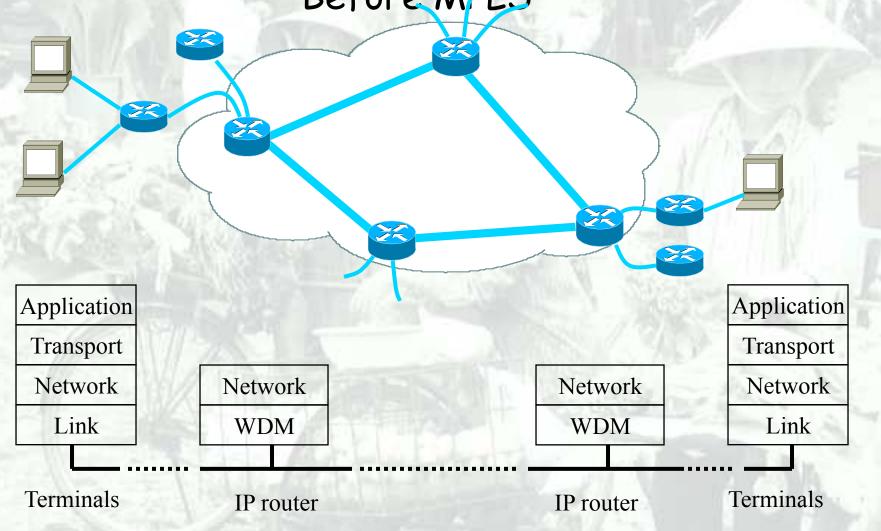
VPN over IP/MPLS

□ IP/MPLS replace dedicated networks

MPLS reduces VPN complexity by reducing routing information needed at provider's



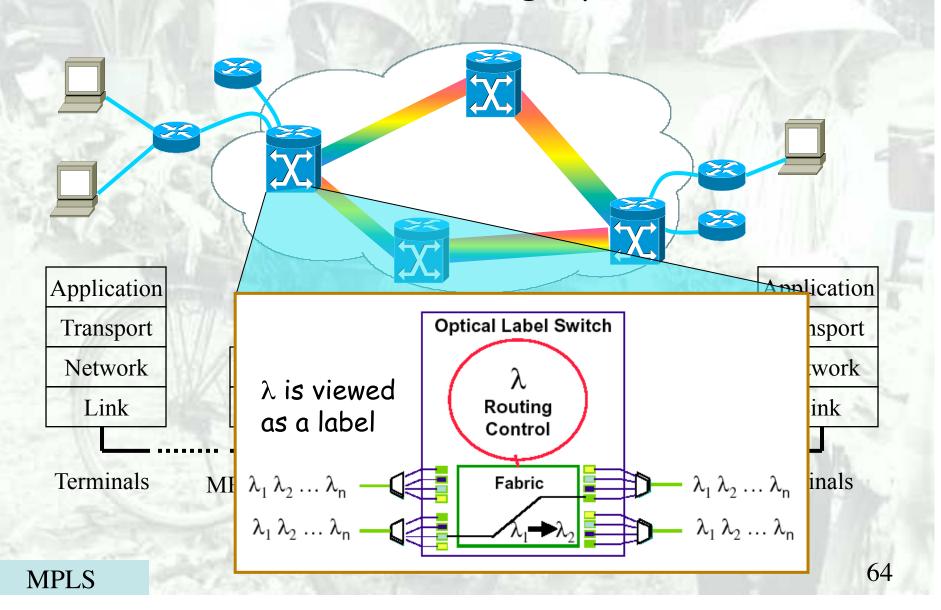
# MPLS for optical networks Before MPLS



Source J. Wang, B. Mukherjee, B. Yoo

### MPLS for ON, con't

MPλS=MPLS+λ lightpath



### MPLS for ON, con't GMPLS

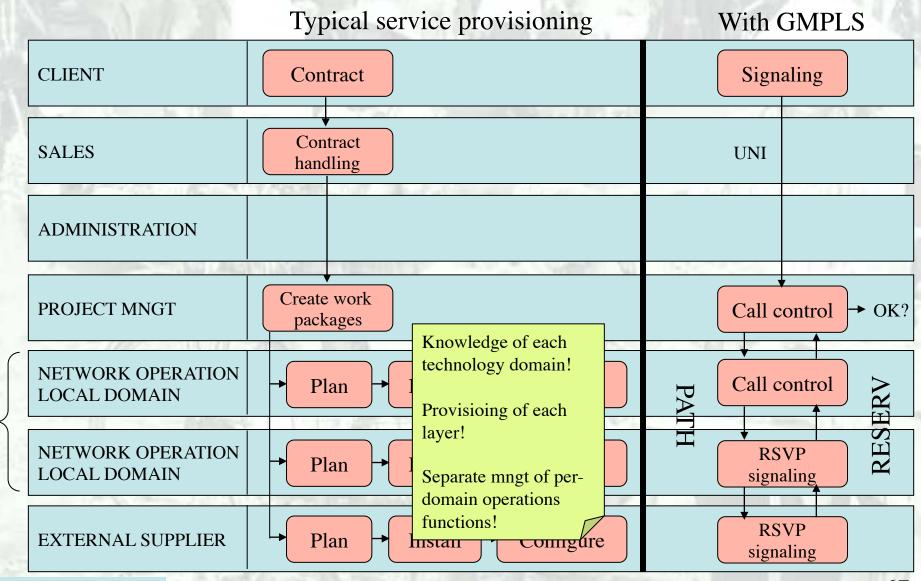
- □ GMPLS stands for "Generalized Multi-Protocol Label Switching"
- Extends the concept of MPLS beyond data networks to address legacy transport networks
- □ Reduce OPEX cost for operators
- ☐ A suite of protocols that provides a common set of control functions for disparate transport technologies (IP, ATM, SONET/SDH, DWDM)
- ☐ Hot issue at IETF!

### MPLS for ON, con't

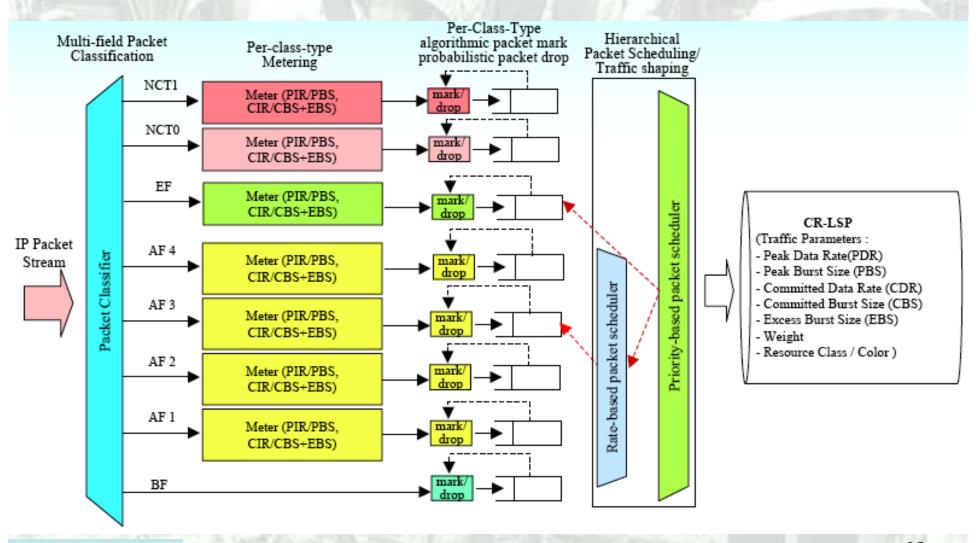
#### GMPLS control plane

LINK MANAGEMENT: Link Management Protocol (LMP)	-Neighbor discovery -Maintain control channel connectivity -Verify data link connectivity -Correlate link property information -Suppress downstream alarms -Localize link failures
ROUTING: Open Shortest Path First-Traffic Engineering (OSPF- TE)	-Distribute TE link information -Advertise nodes in the network and create topology -Calculate constrained shorted path (CSPF) -Routing information for control and data plane
SIGNALING: Resource ReserVation Protocol-Traffic Engineering (RSVP-TE)	-Signals setup/teardown/refresh of paths with QoS requirements (e.g., circuit size) -Uses control channel to setup an optical LSP -Supports refresh reduction -Supports Explicit Route Object (ERO) and Record Route Object (RRO)

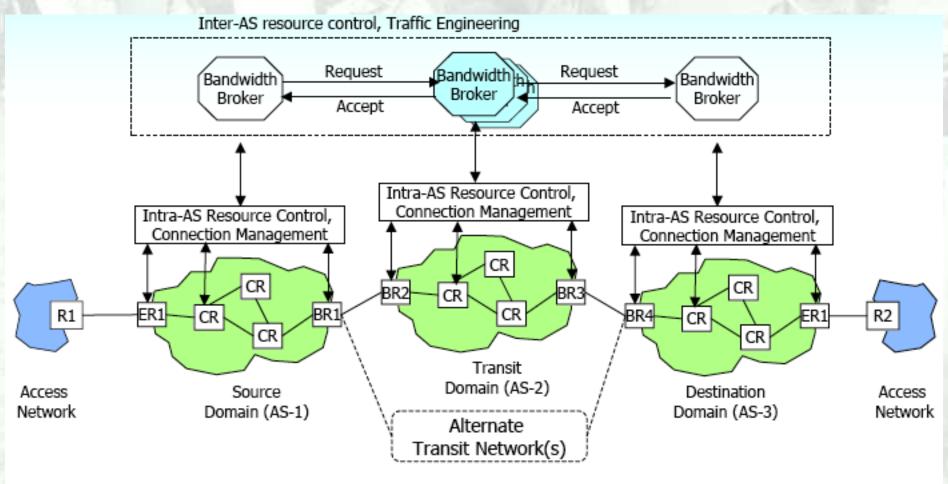
### Ex: Service Provisioning



# DiffServ over (G)MPLS map DiffServ class on MPLS FEC

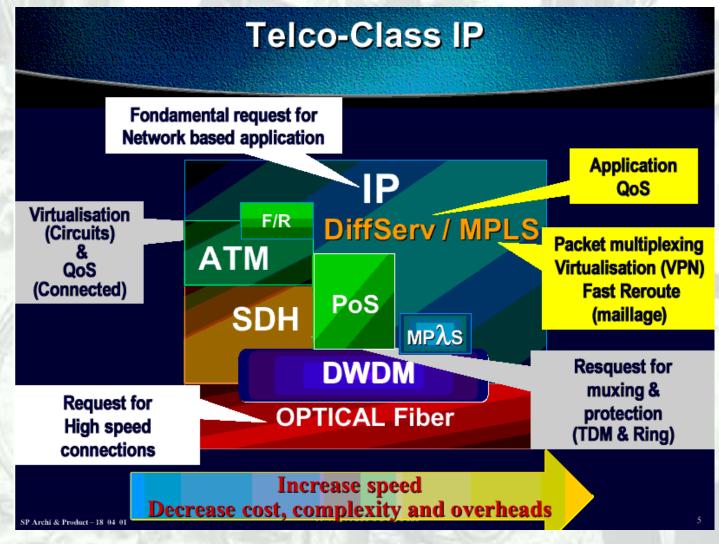


#### Some words on inter-domain

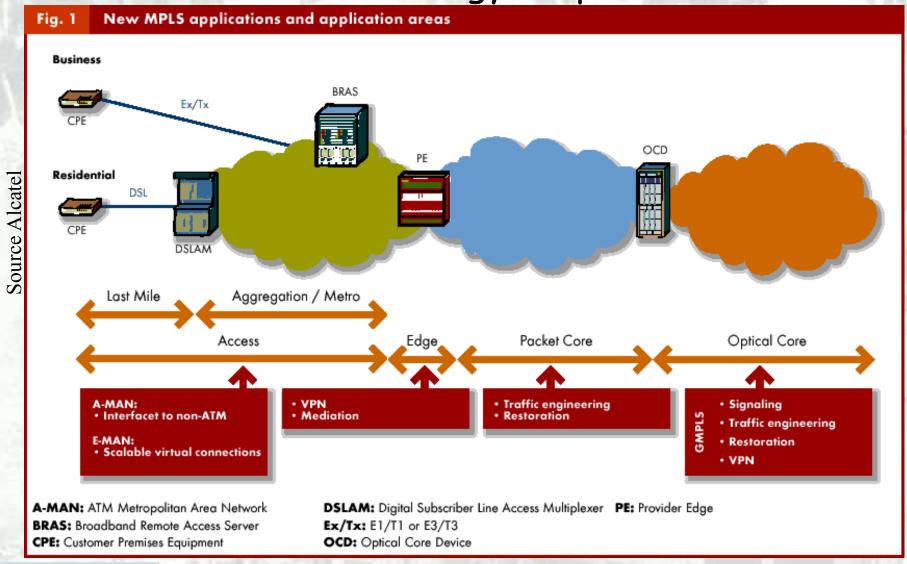


# From cisco

## Summary Towards IP/(G)MPLS/DWDM



#### Summary Technology scope



#### Want to know more?

- □ GMPLS: IEEE Comm. Mag., Vol. 43(7), July 2005
- □ Optical Control Plane for the Grid Community: IEEE Comm. Mag., Vol. 44(3), March 2006.
- "Optical Transport Systems/Networks" by S. Kinoshita & R. Rabbat, APNOMS 2005. http:// www.apnoms.org/2005/tutorial/Tutorial%202.pdf
- « Inter-domain Traffic Engineering for QoSguaranteed DiffServ Provisioning », Young-Tak Kim, APNOMS 2005.
  - http://www.apnoms.org/2005/tutorial/Tutorial %203.pdf
- □ See Tutorial IV of HOTI 2006: Dynamic Optimal Networks for Grid Computing

